

# Dynamic Games with Incomplete Information (continued)

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## Takeover Game

- Target's value  $v$  is uniform on  $[0,100]$
- Target knows value; Acquirer does not
- Acquirer's value is  $1.5 v$
- What offer  $p$  should Acquirer make to Target? This is take-it-or-leave-it bargaining game.

# Auctions

- Auctions are important institution.
- Understanding auctions should help us understand the formation of markets by modeling the competition on one side of the market.
- Auctions represent an excellent application of game theory, since in an auction the rules of the game are made explicit.

## Simple Auctions

Auctions typically take one of four simple forms:

<u>Oral</u>	<u>Sealed Bid</u>
English ( $\uparrow$ price)	2nd Price
Dutch ( $\downarrow$ price)	$\equiv$ 1st Price

# Simple Auctions

- *English*: price increases until only one bidder is left; the remaining bidder gets the good and pays the highest bid.
- *Dutch*: prices decreases until a bidder accepts the price; this bidder gets the good and pays the price at acceptance.
- *Second Price*: each bidder submits a bid in a sealed envelope; the highest bidder gets the good and pays the second highest bid.
- *First Price*: each bidder submits a bid in a sealed envelope; the highest bidder gets the good and pays the amount of his bid.

## Auction Exercise

- Bid for single object
- Common value = \$1 per bean
- On slip of paper write:
  - Name
  - Estimate (# of beans  $\times$  \$1)
  - Bid in first-price sealed-bid auction
  - Bid in second-price sealed-bid auction

# Models of Private Information

- (1) Independent Private Value:  
 $v_i \sim F_i$  independently of  $v_j$  for  $j \neq i$ .
- (2) Common Value:  
 $e_i = v + \varepsilon_i$ ,  $\varepsilon_i \sim F_i$  w/ mean 0.
- (3) Affiliated Value:  
 $v_i(x, s)$ , my value depends on private information  $x = (x_1, \dots, x_n)$  and state of world  $s$ .

## Winner's Curse

I won. Therefore, I overestimated the most. My bid only matters when I win, so I should condition my bid on winning (i.e., that I overestimated the most).

- Winning is bad news about my estimate of value. This is a form of adverse selection that arises in any exchange setting: if you want to trade with me, it must be that no one else offered more, because they did not think that the item is worth what I am willing to pay.

# Models of Private Information

- Independent private value model: It makes sense if differences in value arise from heterogeneous preferences over the attributes of the item
- Common Value: It makes sense if the bidders have homogeneous preferences, so they value the item the same ex post, but have different estimates of this true value.
- Affiliated value model: In this model, each bidder has private information that is positively correlated with the bidder's value of the good.

## Auction Multiple Items

- 2 bidders (L and S), 2 identical items
- L has a value of \$100 for 1 and \$200 for both
- S has a value of \$90 for 1 and \$180 for both
- Uniform-price auction
  - Submit bid for each item
  - Highest 2 bids get items
  - 3<sup>rd</sup> highest bid determines price paid
- Ascending clock auction
  - Price starts at 0 and increases in small increments
  - Bidders express how many they want at current price
  - Bidders can only lower quantity as price rises
  - Auction ends when no excess demand (i.e. just two demanded); winners pay clock price

# What if private information?

- 2 bidders (L and S), 2 identical items
- L has constant marginal value  $u$  drawn  $U[0,1]$
- S has constant marginal value  $v$  drawn  $U[0,1]$
- Uniform-price auction
  - Submit bid for each item
  - Highest 2 bids get items
  - 3<sup>rd</sup> highest bid determines price paid
- Ascending clock auction
  - Price starts at 0 and increases in small increments
  - Bidders express how many they want at current price
  - Bidders can only lower quantity as price rises
  - Auction ends when no excess demand (i.e. just two demanded); winners pay clock price